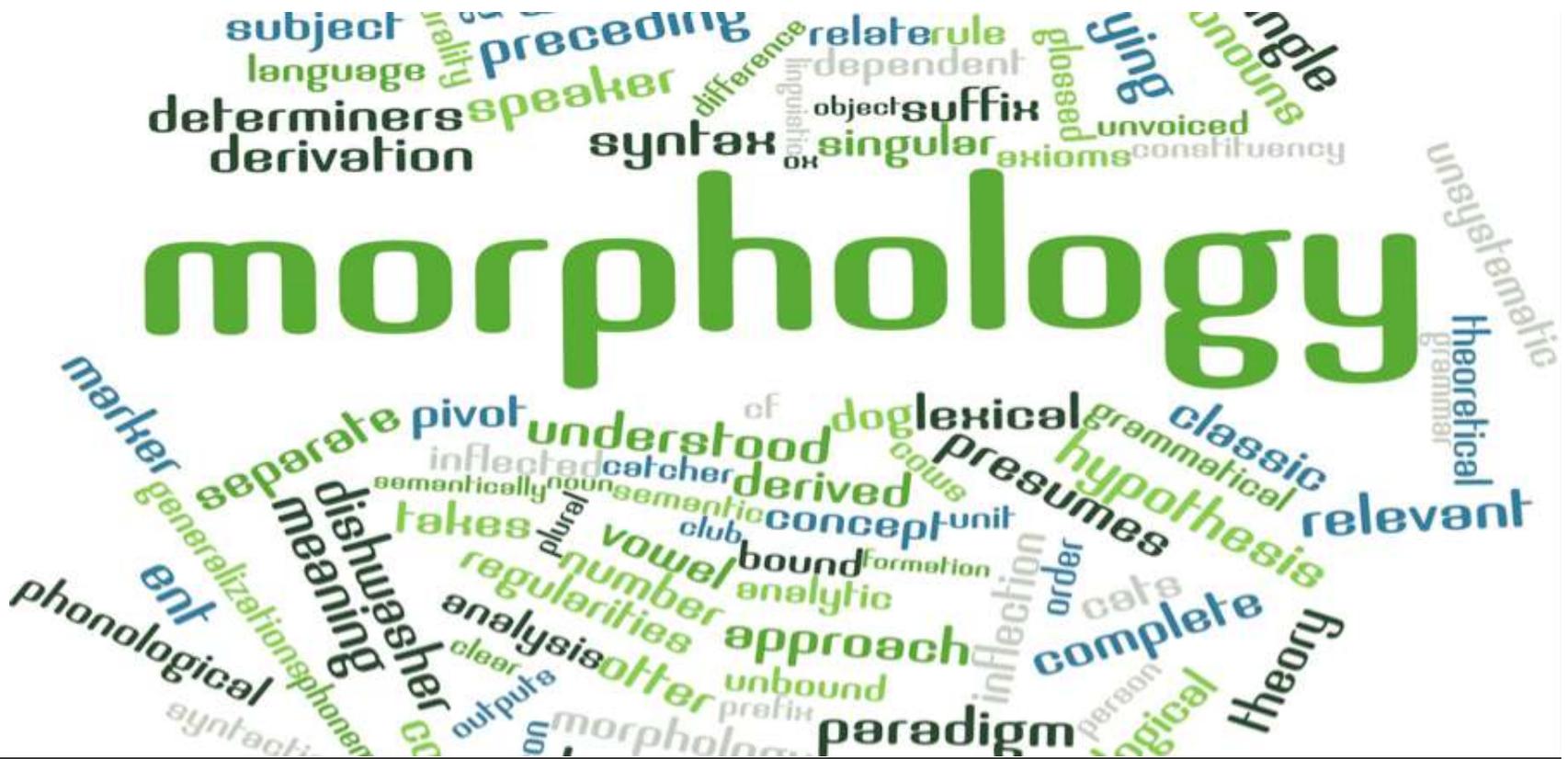


Finite State Morphology

by

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One Specialty: Morphology

Related specialty: phonology

What is Linguistic Morphology?

- Morphology is the study of the internal structure of words.
 - **Derivational morphology.** How new words are created from existing words.
 - [become]
 - [[becom]ing]
 - [un[becom]ing]]
 - **Inflectional morphology.** How features relevant to the syntactic context of a word are marked on that word.
 - **Der Mann** schläft. “The man is sleeping.”
 - Ich sehe **den Mann**. “I see the man.”
 - **Compounding.** Creating new words by combining existing words (not discussed much in this lecture).

Morphemes

- A venerable way of looking at morphology.
- **Morphemes.** Minimal pairings of form and meaning.
 - **Roots.** The “core” of a word that carries its basic meaning.
 - *apple* : ‘apple’
 - *walk* : ‘walk’
 - **Affixes (prefixes, suffixes, infixes, and circumfixes).**
Morphemes that are added to a base (a root or stem) to perform either derivational or inflectional functions.
 - *un-* : ‘NEG’
 - *-s* : ‘PLURAL’

Morphologies of Languages Vary Widely

- Some traditional categories which are problematic but still useful:
 - **Analytic** or **isolating**. Some languages, like [English](#) and especially [Chinese](#), have relatively few affixes (especially inflectional affixes).
 - **Agglutinative**. Most languages (including [Hungarian](#), [Finnish](#), [Turkish](#), [Swahili](#), [Japanese](#), [Korean](#), and most native languages of the Americas, Australia, and New Guinea) perform derivation and inflection by the concatenation of easily-segmentable affixes before, after, and in the root.
 - **Fusional** or **flexional**. Some languages, like [German](#), [Spanish](#), [Russian](#), and [Greek](#), inflect words with the addition of difficult-to-segment *desinences*. For example, the same sequence of segments added to nouns may realize both case and number features.

Isolating Languages: Chinese

- Morphological analyzers and generators are not very useful for isolating languages like English, Vietnamese, Yoruba, and Chinese.
- These languages have little morphology other than compounding.
- **Chinese** has few—if any—affixes (prefixes and suffixes):
 - 们: 我们, 你们, 他们, 。 。 。 同志们
mén: *wǒmén*, *nǐmén*, *tāmén*, *tóngzhìmén*
plural: we, you (pl.), they comrades, LGBT people
 - A couple of “suffixes” that mark aspect: 着 *-zhě* ‘continuous aspect’

Agglutinative Languages: Swahili

Swahili	English
<i>m-tu a-li-lala</i>	'The person slept'
<i>m-tu a-ta-lala</i>	'The person will sleep'
<i>wa-tu wa-li-lala</i>	'The people slept'
<i>wa-tu wa-ta-lala</i>	'The people will sleep'

- Words written without hyphens or spaces between morphemes.
- Orange prefixes mark noun class (like gender, except **Swahili** has nine instead of two or three).
 - Verbs agree with nouns in noun class.
 - Adjectives also agree with nouns.
 - Very helpful in parsing.
- Black prefixes indicate tense.

Fusional Languages: A New World Spanish

	Singular			Plural		
	1 st	2 nd	3 rd formal 2 nd	1 st	2 nd	3 rd
Present	<i>am-o</i>	<i>am-as</i>	<i>am-a</i>	<i>am-a-mos</i>	<i>am-áis</i>	<i>am-an</i>
Imperfect	<i>am-ab-a</i>	<i>am-ab-as</i>	<i>am-ab-a</i>	<i>am-áb-a-mos</i>	<i>am-ab-ais</i>	<i>am-ab-an</i>
Preterit	<i>am-é</i>	<i>am-aste</i>	<i>am-ó</i>	<i>am-a-mos</i>	<i>am-asteis</i>	<i>am-aron</i>
Future	<i>am-aré</i>	<i>am-arás</i>	<i>am-ará</i>	<i>am-are-mos</i>	<i>am-aréis</i>	<i>am-arán</i>
Conditional	<i>am-aría</i>	<i>am-arías</i>	<i>am-aría</i>	<i>am-aría-mos</i>	<i>am-aríais</i>	<i>am-arían</i>

Root-and-Pattern Morphology

- **Root-and-pattern.** A special kind of fusional morphology found in Arabic, Hebrew, and their cousins.
- Root usually consists of a sequence of consonants.
- Words are derived and, to some extent, inflected by patterns of vowels intercalated among the root consonants.
 - **kitaab** 'book'
 - **kaatib** 'writer; writing'
 - **maktab** 'office; desk'
 - **maktaba** 'library'

Other Non-Concatenative Morphological Processes

- Root-and-pattern morphology is an extreme case of the general class of morphology that is most difficult to model using computational methods.
- **Non-concatenative morphology** involves operations other than the concatenation of affixes with bases.
 - **Infixation**.
 - **Reduplication**. Can be prefixing, suffixing, and even infixing.
 - **Internal change** (tone change; stress shift; apophony, such as umlaut and ablaut).
 - **Root-and-pattern** morphology.
 - And more...

Polysynthetic Morphologies

- Polysynthetic morphologies are an interesting special case of agglutinative languages.
- Not only do they allow the formation and inflection of words with multiple, segmentable morphemes, they allow the creation of full “sentences” by morphological means.
- They often allow the incorporation of nouns into verbs.
- They may also have affixes that attach to verbs and take the place of nouns.
- **Yupik Eskimo**
untu-ssur-qatar-ni-ksaite-ngqiggte-uq
reindeer-hunt-FUT-say-NEG-again-3SG.INDIC
'He had not yet said again that he was going to hunt reindeer.'

Language Comparison wrt FSTs

- ❑ Morphologies of all types can be analyzed using finite state methods.
- ❑ Some present more challenges than others.
 - ❑ **Analytic languages.** Trivial, since there is little or no morphology (other than compounding).
 - ❑ **Agglutinating languages.** Straightforward—finite state morphology was “made” for languages like this.
 - ❑ **Polysynthetic languages.** Similar to agglutinating languages, but with blurred lines between morphology and syntax.
 - ❑ **Fusional languages.** Easy enough to analyze using finite state method as long as one allows “morphemes” to have lots of simultaneous meanings and one is willing to employ some additional tricks.
 - ❑ **Root-and-pattern languages.** Require some very clever tricks.

Why Compute When Storage is Cheap?



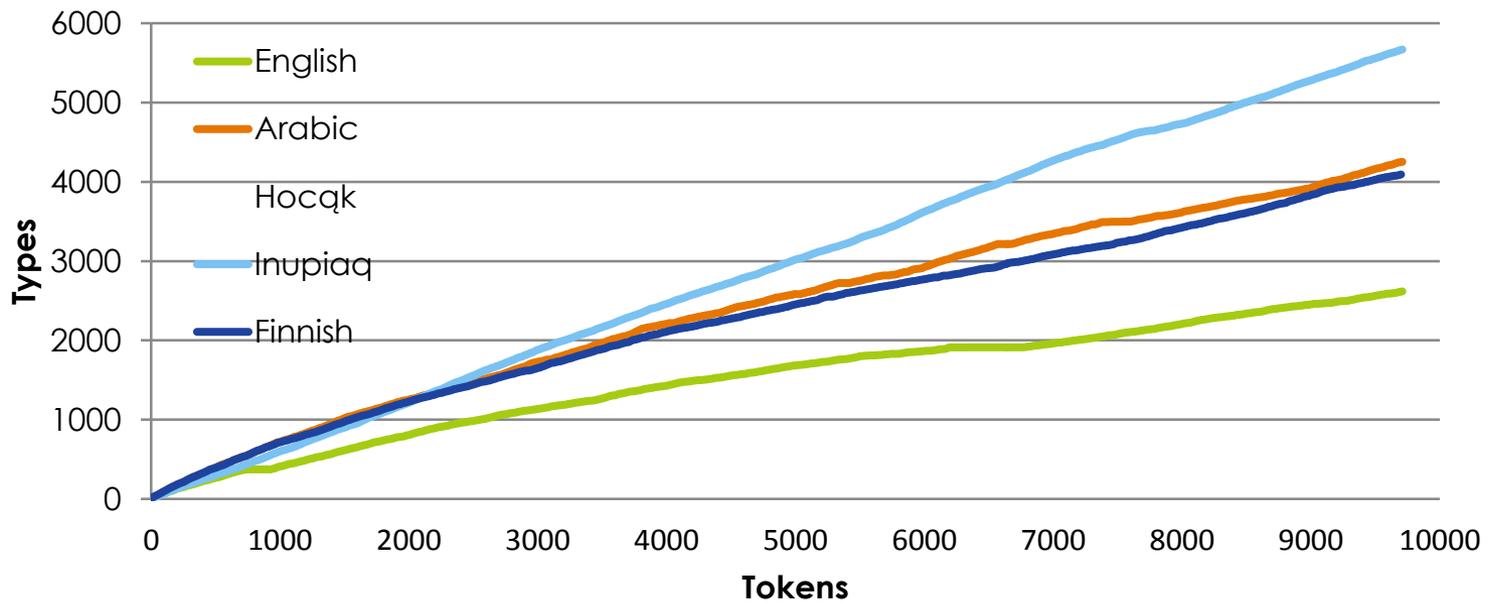
English is Easy

- You could get by performing significant tasks in English without implementing much English morphology.
- You can store most of the derived forms of words in memory instead of their inflected forms.
- English **does** have **productive morphological constructions**, so you could encounter some out of dictionary issues, but for the most part, things would come out in the wash.
- Chinese is even easier.
- **However, there are other languages.**

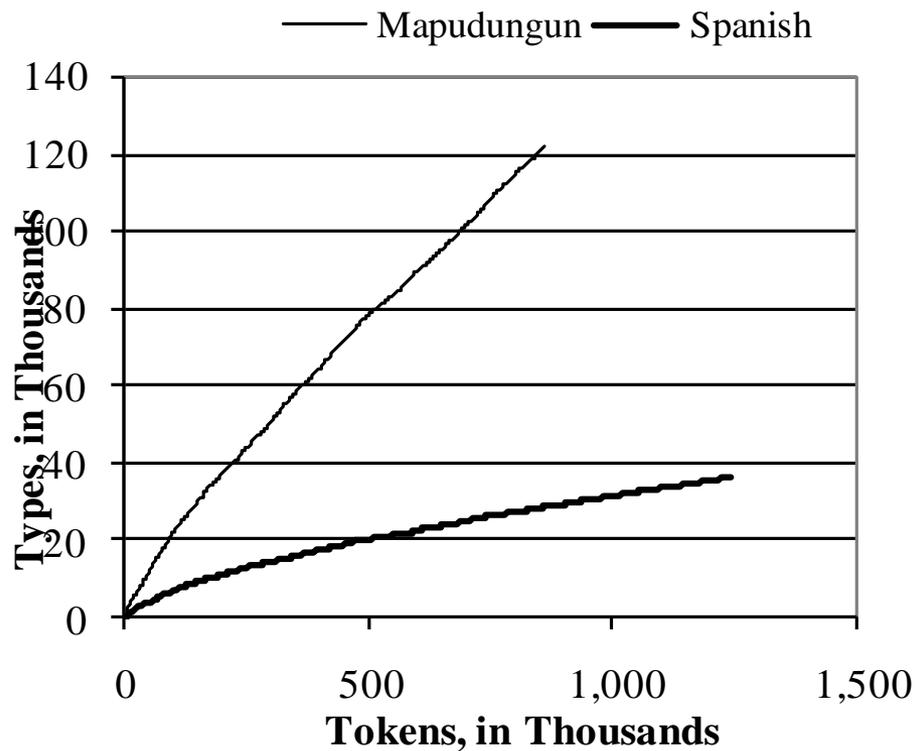
Morphological patterns that can be employed to derive or inflect new words.

Inupiaq Needs an Analyzer/Generator

Type-Token Curves



Mapathungun Is No Easier



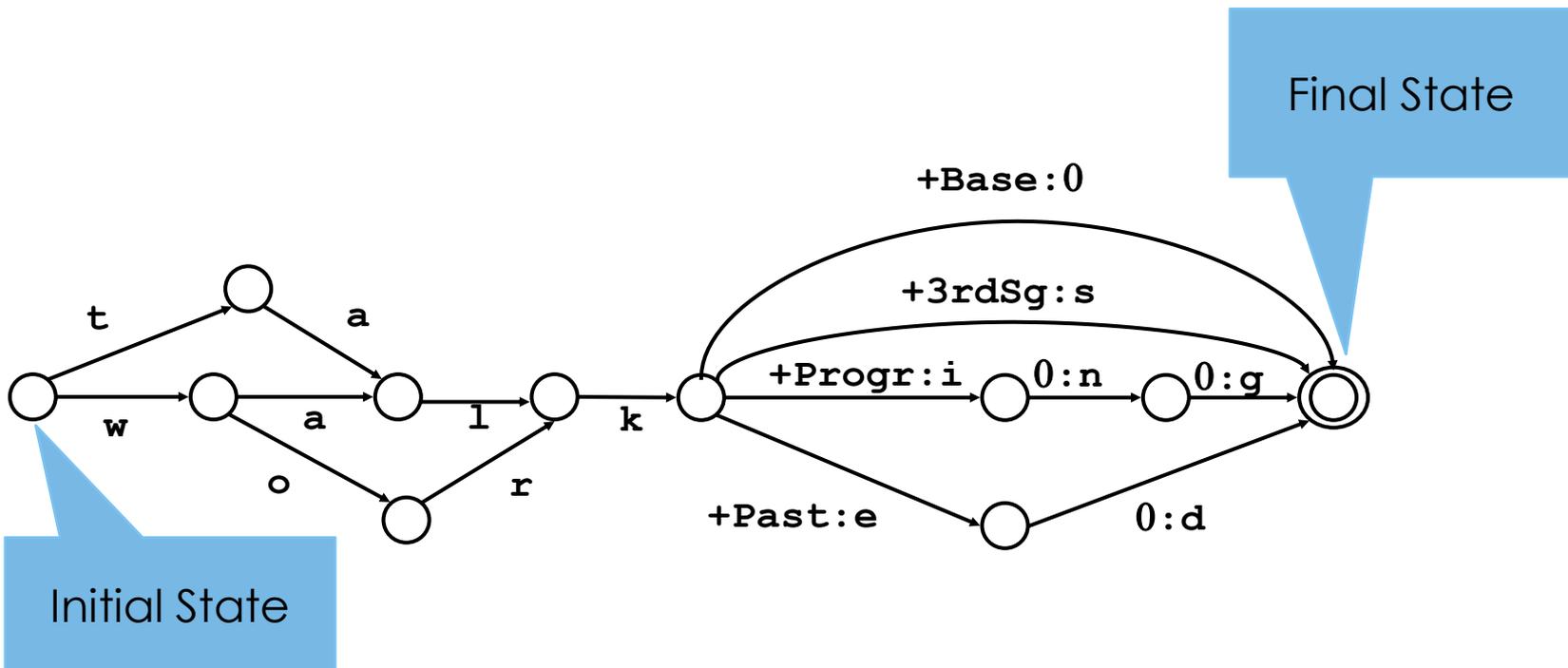
The Good News

- More than almost any other problem in computational linguistics, morphology is a solved problem (as long as you can afford to write rules by hand).
- Finite state morphology is one of the great successes of natural language processing.
- One brilliant aspect of using FSTs for morphology: the **same code** can handle both **analysis** and **generation**.

Finite State Morphology and Finite State POS Tagging

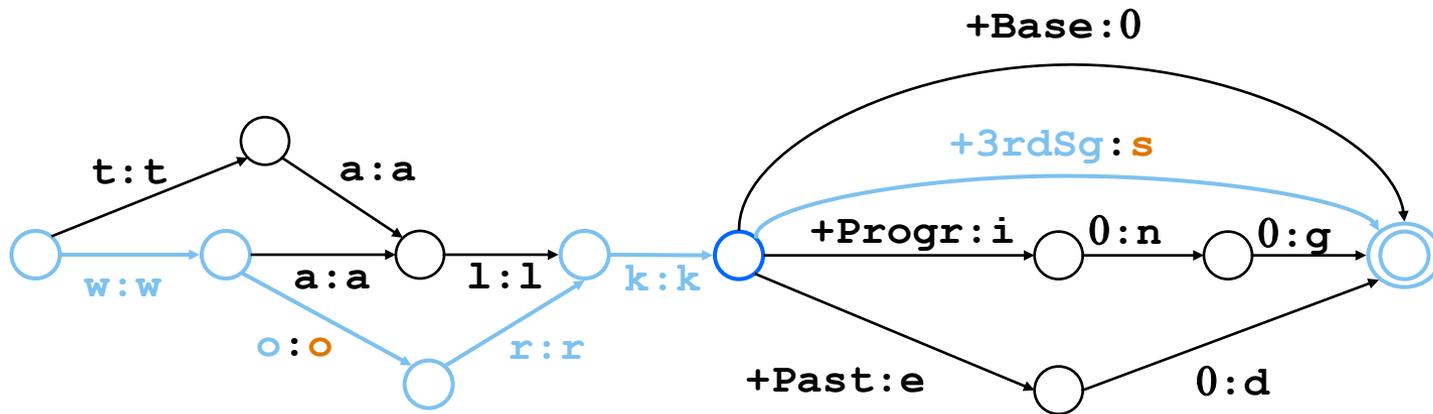
- ❑ **Morphological analysis** is not the same as **part of speech tagging**, though they are related in some ways.
- ❑ **Part of speech tagging**.
 - ❑ **Tokenization**. Break text into tokens.
 - ❑ **Lexical lookup**. Lookup each tag for each token.
 - ❑ **Disambiguation**. Choose tag for token, often on basis of statistical model (e.g. HMM).
- ❑ **Morphological analysis**.
 - ❑ Transduce each token as stem/translation and a sequence of morphological properties.
 - ❑ **amaríamos** → **love+Conditional+1P+Plural**

A Finite State Transducer

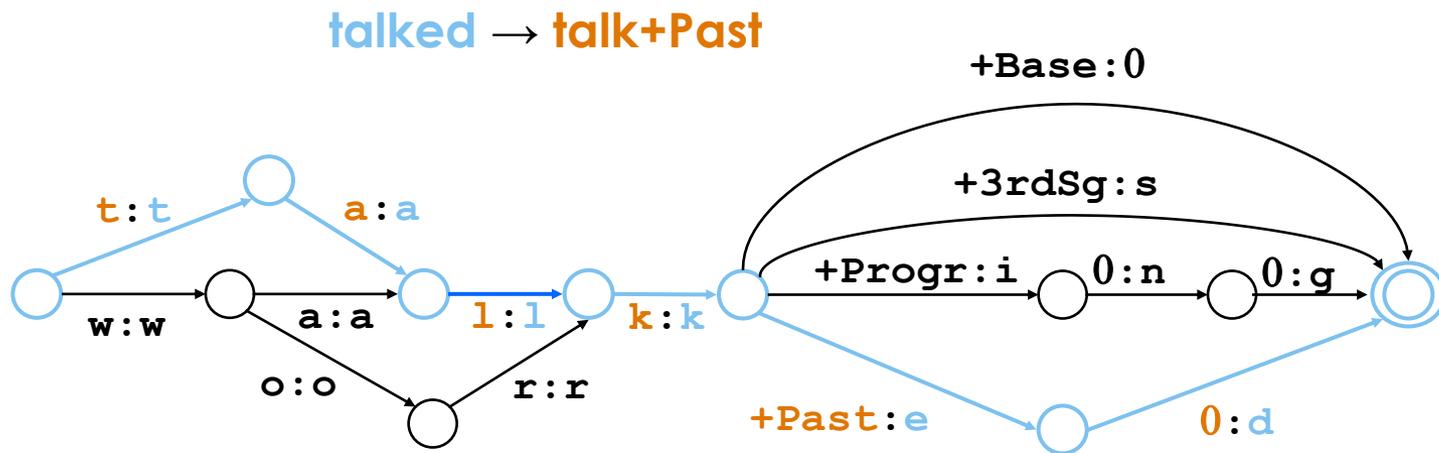


Generation

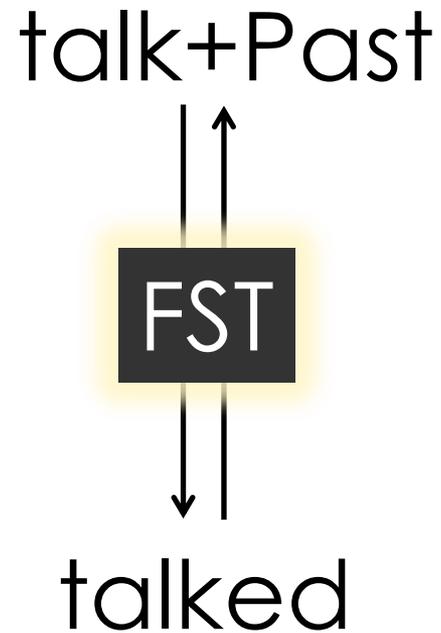
work+3rdSg → works

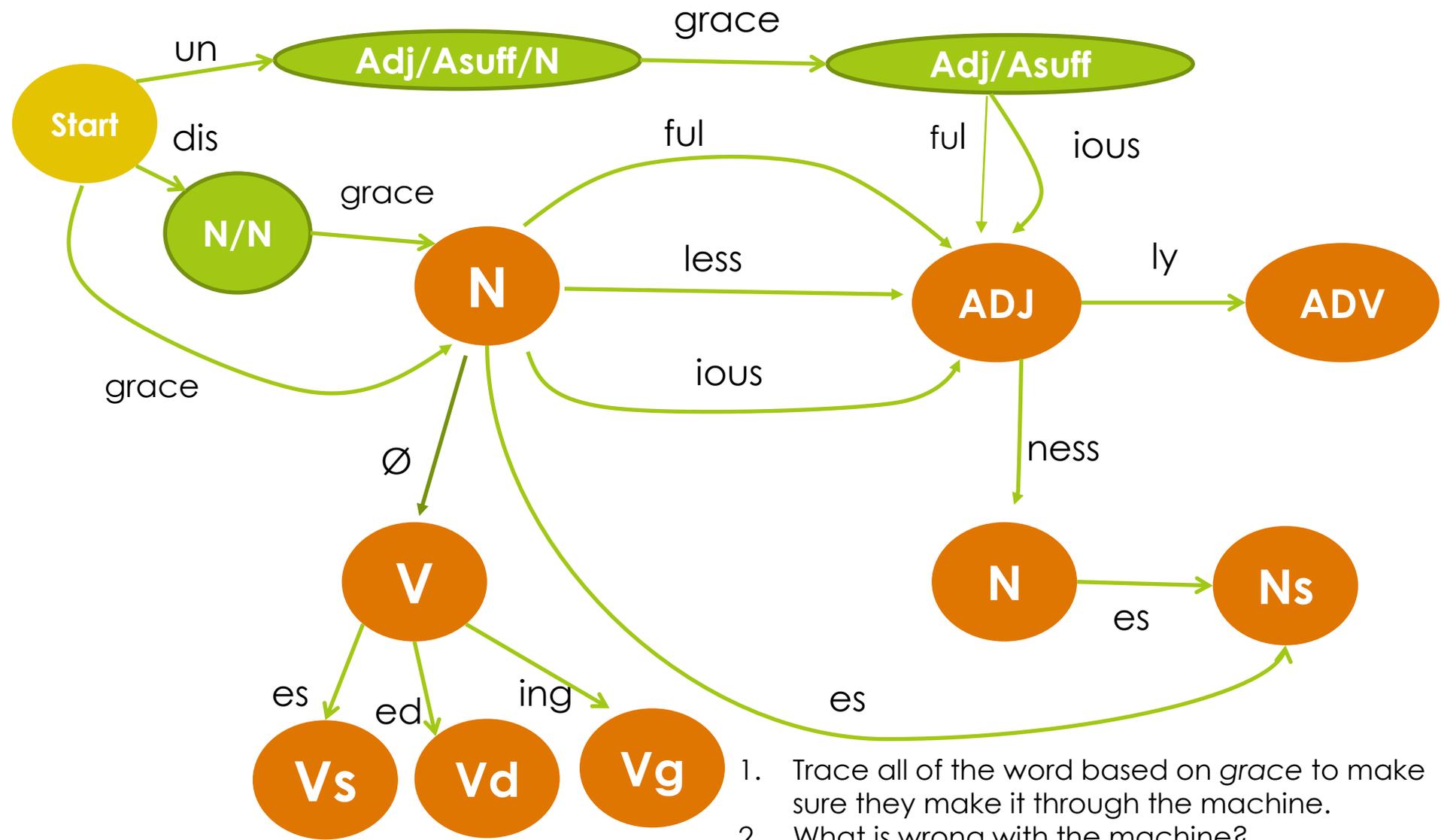


Analysis



Upper Side/Lower Side





1. Trace all of the word based on *grace* to make sure they make it through the machine.
2. What is wrong with the machine? Overgeneration? Undergeneration? Redundancy? Does not capture structure and ambiguity?

Tools

- There are special finite state toolkits for building morphological tools (and other linguistic tools).
- The best-known of these is the **Xerox Finite State Tool** or **XFST**, which originated at Xerox PARC.
- There are open source reimplementations of XFST called **HFST** (Helsinki Finite State Technology) and **Foma**, which are not as fully optimized as XFST but which are sometimes more pleasant to use.
- None of these tools allow the construction of weighted FSTs.

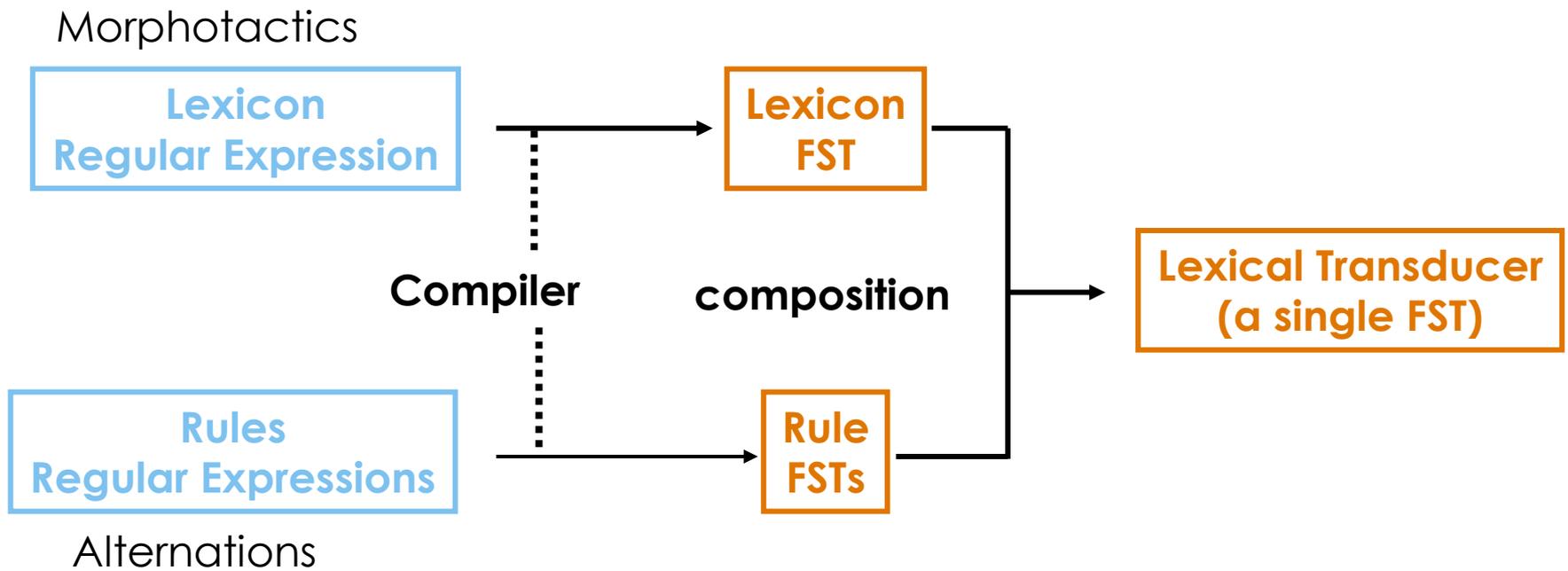
A Convenient Division I

- XFST makes a useful division:
 - **lexc**: constructs finite state transducers from lexical rules defining morphemes and how they can be concatenated.
 - **xfst**: constructs finite state transducers that implement phonological alternations (changes in the segments—vowels and consonants—in a word that come about when morphological operations change their context). The most common phonological alternations occur at morpheme boundaries; however, phonological alternations can also be non-local.
 - /san + pat/ → [səmpət]
 - /dæg + z/ → [dægz], /kæt + z/ → [kæts], /haʊs + z/ → [haʊsəz]

A Convenient Division II

- The FSTs produced by **lexc** and **xfst** have the same data formats and can be straightforwardly composed, etc.
- **lexc** is designed to define lexicons and the relationship between **morphemes** and their **properties**.
- FSTs can do most, but not quite all, of the things the phonological rules traditionally used by linguists can do.
Specialty of xfst.
- **This distinction may be useful in future tools for finite state morphology** even though it is not algorithmic in nature.

How Lexical Transducers are Made



Phonological Rule Ordering

- Traditional phonological rules, in a formalism proposed by Morris Halle and Noam Chomsky, are ordered rewrite rules.
- $\mathbf{a} \rightarrow \mathbf{b} / \mathbf{X_Y}$ (where \mathbf{a} is rewritten as \mathbf{b} between strings \mathbf{X} and \mathbf{Y})
- **Orderings and interactions:**
 - The way rules are ordered relative to other interacting rules produces interesting patterns.
 - Rules can also interact with themselves, assuming that they apply either **RTL** or **LTR** instead of **simultaneously**.

Kinds of Interactions I

- **Feeding**

$Rule_i$ creates environments in which $Rule_{i+n}$ can apply.

- **Bleeding**

$Rule_i$ destroys environments in which $Rule_{i+n}$ could apply.

- These interactions are very common in natural language.

Kinds of Interactions II

□ Counter-feeding

$Rule_{i+n}$ would create environments in which $Rule_i$ would apply if $Rule_{i+n}$ were ordered before $Rule_i$.

□ Counter-bleeding

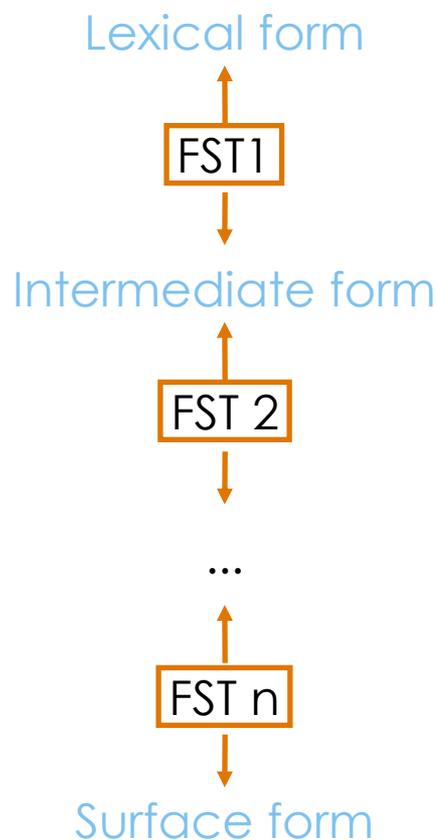
$Rule_{i+n}$ would destroy environments in which $Rule_i$ would apply if $Rule_{i+n}$ were ordered before $Rule_i$.

- What kind(s) of interaction that is/are possible with the Chomsky and Halle formalism can't be modeled via the composition FSTs?

Self-Interactions

- Interactions between two or more rules are always possible via composition of FSTs.
- However, self-interactions where a rule feeds or bleeds itself cannot be implemented with FSTs.
- Fortunately, relatively few phonological alternations are implemented using self-feeding (and even fewer, using self-bleeding).
 - **Spreading vowel harmony.** When vowels change to be more like neighboring vowels. This can apply iteratively.
 - **Tone spreading.** When tones change to match neighboring tones. This can also apply iteratively.

Ordered Rules Modeled as Cascade of FSTs



Ordered sequence of rewrite rules (Chomsky & Halle 1968) can be modeled by a cascade of finite-state transducers (Johnson 1972; Kaplan & Kay 1981)

Sequential Application

k a N p a n



k a m p a n



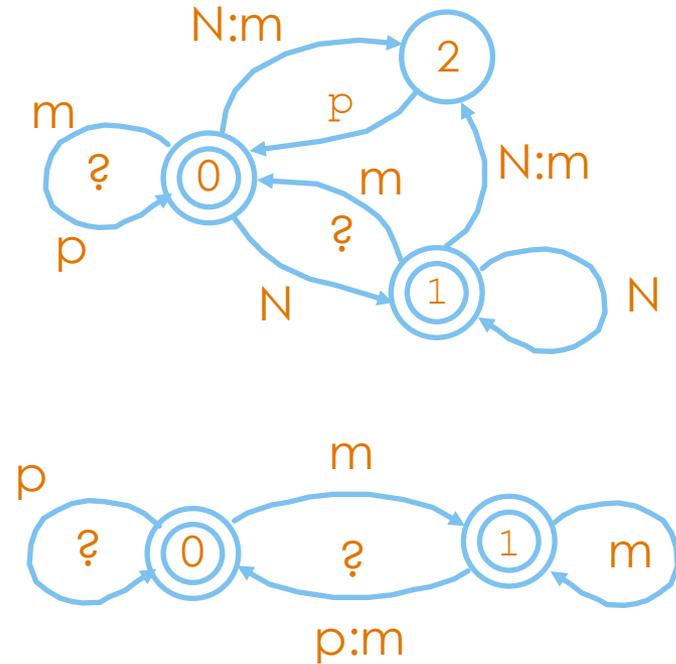
k a m m a n

$N \rightarrow m / _ p$

$p \rightarrow m / m _$

Sequential Application in Detail

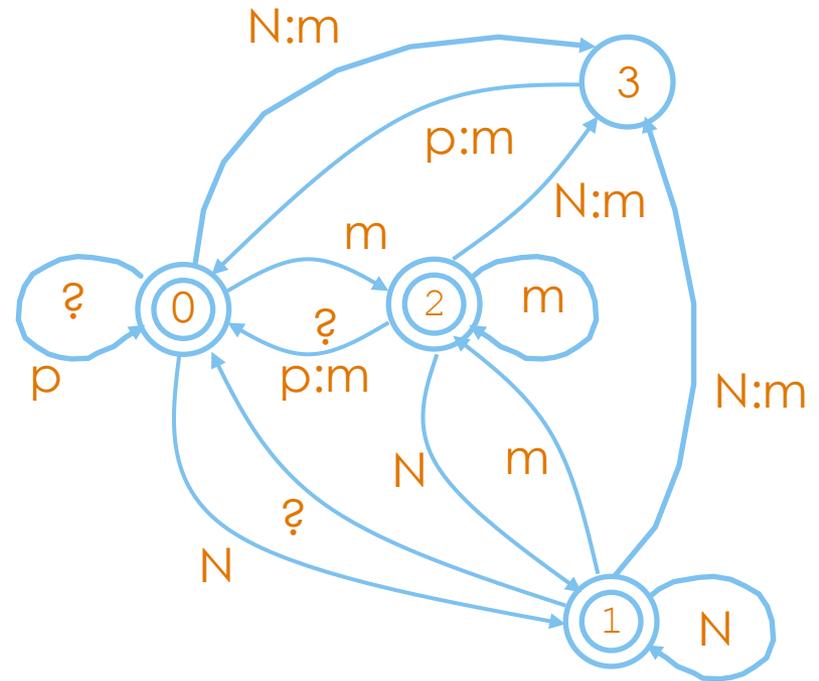
k a N p a n
 0 0 0 | 2 0 0 0
 ↓
 k a m p a n
 0 0 0 1 | 0 0 0
 ↓
 k a m m a n



Composition

k a N p a n
0 0 0 3 0 0 0
k a m m a n

Diagram illustrating the composition of the word "kammann" from the sequence "k a N p a n". The sequence "k a N p a n" is shown above "0 0 0 3 0 0 0", which is aligned with "k a m m a n". Vertical arrows point from the 'a' and 'N' in the top sequence to the 'a' and 'm' in the bottom sequence, respectively.



A Cheat: Flag Diacritics

- Sometimes, it is easier to write FSTs for morphology when the FST formalism is enhanced with **flag diacritics**.
- These essentially add memory to finite state machines.
- They are allowed in **XFST**, **Foma**, and **HSFT**. Elsewhere, too.
 - If you look at XFST (**lexc**) scripts, you will see strings like **@U.CASE.ACC@**, with alphanumeric characters and full stops enclosed in “@” characters.
 - This is a unification flag; there are other types of flags.
 - “Traversing an arc labeled with a Flag Diacritic is like an epsilon transition but is conditioned on the success or failure of an OPERATION specified by the Flag Diacritic. The result depends on the state of a **feature register** that is initialized in the beginning of the analysis or generation and is contiguously updated along each path that is being explored.” (Beesley and Karttunen, 2003)

Flag Diacritic Example: Arabic

'book'	analysis
<i>kitaab+u</i>	definite nominative
<i>kitaab+a</i>	definite accusative
<i>kitaab+i</i>	definite genitive
<i>kitaab+uN</i>	indefinite nominative
<i>kitaab+aN</i>	indefinite accusative
<i>kitaab+iN</i>	indefinite genitive
<i>al+kitaab+u</i>	definite prefix + definite nominative
<i>al+kitaab+a</i>	definite prefix + definite accusative
<i>al+kitaab+i</i>	definite prefix + definite genitive

A Constraint

- The definite prefix *al-* cannot co-occur with the indefinite case markers *-uN*, *-an*, and *-in*.
- It is possible to capture this distribution with pure finite state technology.
- Not terribly fun.
- Instead...
 - Set a feature register using a flag diacritic `@U.ART.PRESENT@` when you consume the definite prefix .
 - Insure that indefinite case suffixes cannot be present if the feature register is set accordingly.

Arabic Definiteness in XFST's lexc I

Multichar_Symbols @U.ART.PRESENTS@ @U.ART.ABSENT@ uN aN
iN

LEXICON Root
Article ;

LEXICON Article
al@U.ART.PRESENT@ Stem ; ! optional article prefix
Stem ; ! empty string entry

LEXICON Stems
kitaab Case ; ! one stem to represent > 10,000

Arabic Definiteness in XFST's lexc II

LEXICON Case

u # ;

a # ;

i # ;

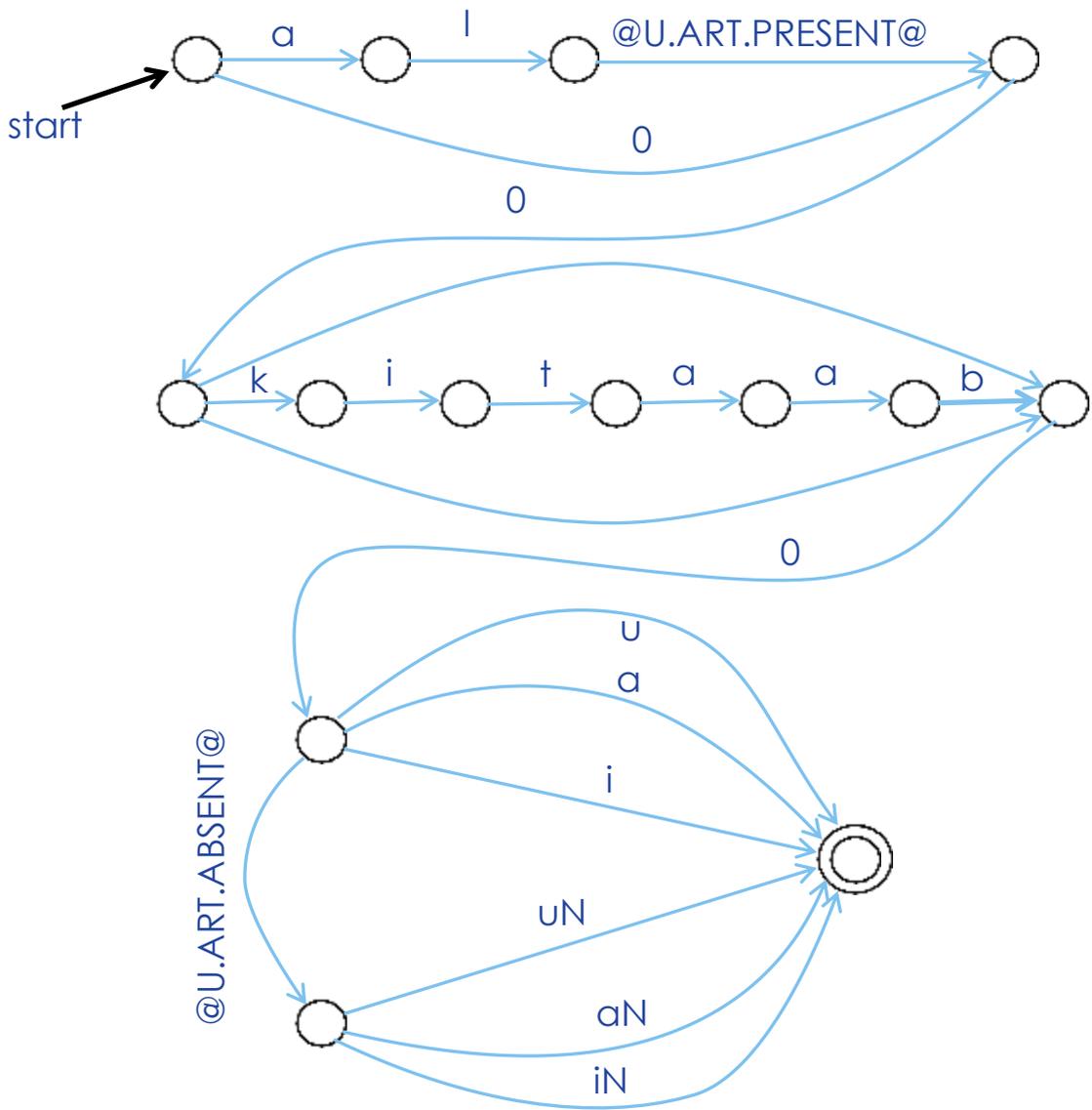
@U.ART.ABSENT@ IndefCase ;

LEXICON IndefCase

uN # ;

aN # ;

iN # ;



Arabic FSM Fragment

Upshot

While **flag diacritics** are not a fundamental part of the finite state paradigm, they are a useful addition in cases where humans will manually code finite state morphological analyzers or generators. They will be used in our finite state implementations of non-concatenative morphology.



Concatenation is Regular

- The most common kind of morphological operation involves the concatenation of morphemes.
 - Agglutinative and polysynthetic languages use concatenation primarily or exclusively.
 - Concatenation is straightforwardly regular.
 - We can easily implement it with finite state methods.

More Than Mere Concatenation

- **Reduplication.** Repeating all or part of a word as a morphological operation.
- **Infixation.** Inserting an **affix** into a **base**.
- **Root-and-pattern morphology.** Interdigitation in Semitic and its relatives.
- Others. **Apophony**, including the umlaut in English *tooth* → *teeth*; **subtractive morphology**, including the **truncation** in English nickname formation (*David* → *Dave*); and so on.

The Problem

- Non-concatenative morphology cannot be addressed with finite state methods without considerable ingenuity.
- However, many languages feature some non-concatenative processes.
 - **Reduplication**—extremely widespread.
 - **Root-and-pattern**—occurs in two very important languages.

A Solution

- Antworth (1990) found a way of capturing reduplication as long as the reduplicant always had the same number of segments (vowels and consonants). **Unsatisfying.**
- Beesley and Karttunen (2000, 2003) proposed the **compile-replace** algorithm, a clever solution that allows purely finite state methods to deal with a wide class of non-concatenative morphology.
- **Summary of compile-replace:**
 - Regex replacements generate new regex, which encodes meta-morphotactic information.
 - New regex is compiled to FSA.
 - FSA describes complete lower side.

Concatenative Morphology Using XSFT: Simplest Case

Initial lexicon (from *lexc*) with unfinished lower side.

.o.

Alternation rules mapping from the lower side of the lexicon FST to surface strings.

Concatenative Morphology Using XFST: Realistic Case

Rules and filters to
remove overgeneration.

.o.

Initial lexicon (from [lexc](#))
with unfinished lower
side.

.o.

Alternation rules
mapping from the lower
side of the lexicon FST to
the surface strings.

Adding Compile-Replace to XSFT

Rules and filters to remove overgeneration.

.o.

Initial lexicon (from `lexc`) with unfinished lower side.

Application of **compile-replace** to resolve meta-morphotactic descriptions on the lower side of the lexicon FST.

.o.

Alternation rules mapping from the lower side of the lexicon FST to surface strings.

Total Reduplication in Malay

Root	Gloss	Reduplication	Gloss
<i>anak</i>	'child'	<i>anakanak</i>	'children'
<i>lembu</i>	'cow'	<i>lembulembu</i>	'cows'
<i>buku</i>	'book'	<i>bukubuku</i>	'books'
<i>basikal</i>	'bicycle'	<i>basikalbasikal</i>	'bicycles'

Script in lexc for Malay Reduplication

Multichar_Symbols **^[^]** +Noun +Unmarked +Plural

LEXICON Root

0:**^[**{ NRoots ;
NRoots ;

LEXICON NRoots

buku NSuff ;
lembu NSuff ;
pelabuhan NSuff ;

LEXICON NSuff

+Noun:0 Num ;

LEXICON Num

+Unmarked:0 # ;
+Plural:**}^2^]** # ;

In retrospect, that
was not a good idea

Script

Though when all the
morphological and
phonological alternations

in Malayalam
incorporate
similar

This is not trivial but is
nevertheless left as an
exercise to the viewer.
See, also, Beesley and
Karttunen (2003).

Upper
Lower

Upper
Lower

Upper
Lower:

But What about Arabic?

Pattern:

a a

Template:

C V C V C

Root:

How would you model this kind of pattern using the **compile-replace** algorithm (or some other finite state method)?

One Solution (Beesley & Karttunen, 2003)

- Start with a network that constructs regexes that, when compiled, “merges” roots and patterns into templates.
- $^{\wedge}\{[ktb].m>.\{CVCVC\}.<m.[a+]^{\wedge}\}$
 - $.m>$. rightward merge
 - $.<m$. leftward merge
- Only works if C and V are defined as “super-symbols”:
 - list C b t y k l m n f w r z d s
 - list V a i u
- $.m>$. “spreads” k, t, and b LTR across the Cs (which match).
- $.<m$. “spreads” a RTL across the Vs (which match).

More Compile-Reduce

- ▣ Just as with Malay total reduplication, **compile-replace** comes to the rescue.
 - ▣ Network generates a regular expression as it consumes the upper side.
 - ▣ **Compile-replace** compiles all (or part) of this expression.
 - ▣ The resulting FSA defines a language consisting of one word.

I Like to Merge It, Merge It

Multichar_Symbols ^[^] +3P +Masc +Fem +Sg +Act +Pass +FormI +FormIII

LEXICON Root

LBound ;

LEXICON LBound

[:^]{ Roots ;

LEXICON Roots

ktb MergeRight ;

drs MergeRight ;

LEXICON MergeRight

0:}.m%>.{ Template ; ! % literalizes the >

I Like to Merge It, Merge It

LEXICON Template

```
+FormI:CVCVC MergeLeft ;  
+FormII:CVVCVC MergeLeft ;
```

LEXICON MergeLeft

```
0:}.%<m. Vocalization ; ! % literalizes the <
```

LEXICON Vocalization

```
+Act:[a+] RBound ;  
+Pass:[u*i] RBound ;
```

LEXICON RBound

```
]:^] PerfEnd ;
```

LEXICON PerfEnd

```
+3P+Masc+Sg:a # ;  
+3P+Fem+Sg:at # ;
```

Conclusion

- Finite state methods provide a simple and powerful means of generating and analyzing words (as well as the phonological alternations that accompany word formation/inflection).
- Straightforward concatenative morphology is easy to implement using finite state methods.
- Other phenomena are easiest to capture with extensions to the finite state paradigm.
 - Co-occurrence restrictions—**flag diacritics**.
 - Non-concatenative morphology—**compile-replace** algorithm. Pure finite state, but computed in a novel fashion.